# Difference-isomorphic graph families 

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Joint work with Lior Gishboliner and Benny Sudakov
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## Set families with certain properties

## Theorem (Sperner 1928)

If $\mathcal{A}$ is a family of distinct subsets of $[n]$ s.t. no set is contained in the other, then $|\mathcal{A}| \leq\binom{ n}{\lfloor/ 2\rfloor}$.

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## Theorem (Erdős-Ko-Rado 1938)

Let $n \geq 2 r$. If $\mathcal{A}$ is a family of distinct $r$-element subsets of $[n]$ s.t. each two subsets intersect, then $|\mathcal{A}| \leq\binom{ n-1}{r-1}$.

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- If $|\mathcal{A}| \approx\binom{n-1}{r-1}$, then $\mathcal{A}$ is "close" to the extremal example.


## Graph families with intersection properties

## Conjecture (Simonovits-Sós 1976)

If $\mathcal{G}$ is a family of graphs on [n] s.t. any two graphs in $\mathcal{G}$ share a common triangle $(\triangle)$, then $|\mathcal{G}| \leq 2^{\binom{n}{2}-3}$.

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- $\triangle \Rightarrow C_{n}, \quad|G| \leq 2^{\binom{n}{2}-n}$ by Leader, Ranđelović and Tan.


## Families with difference properties

## Conjecture (Gowers, 2009)

$\forall \delta>0 \forall n \gg_{\delta} 1$, if $\mathcal{A}$ is a family of $n \times n$ matrices with entries in $\{1, \ldots, k\}$ with $|\mathcal{A}|>\delta \cdot k^{n^{2}}$, then $\exists A_{1}, \ldots, A_{k} \in \mathcal{A}, X \subseteq[n]$ s.t. $A_{i+1}-A_{i}=\mathbb{1}_{X \times X}$ for all $i$.

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## Graph families with symmetric difference properties

Notation: $G_{1} \oplus G_{2}:=\left([n],\left(E\left(G_{1}\right) \backslash E\left(G_{2}\right)\right) \cup\left(E\left(G_{2}\right) \backslash E\left(G_{1}\right)\right)\right)$.

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## Question(Alon 2023+)

When $\mathcal{H}$ contains all the $K_{4}$ 's, does $|\mathcal{G}|=o\left(2{ }_{\binom{n}{2}}^{\text {' }}\right)$ ?

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- If so, maybe $|\mathcal{G}| \leq n^{O(n)}$ as there are $n$ ! isomorphisms.


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- All the above can be extended to families of $r$-graphs $(n \gg r)$.


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- $e_{\psi}\left(N_{\varphi}(G)\right):=\#\left\{G_{1}, G_{2} \in \mathcal{G}: G \stackrel{\varphi}{\cong} G_{1}, G \stackrel{\varphi}{\cong} G_{1}, G_{1} \stackrel{\psi}{\cong} G_{2}\right\}$.
- Let $M=2^{\left.\frac{1}{2}\binom{n}{2}-\left\lfloor\frac{n}{2}\right\rfloor\right)}$ be the extremal number.


## Lemma (technical)

- $e_{\psi}\left(N_{\varphi}(G)\right)<M^{2} \cdot n^{-10 n}$ unless $\varphi, \psi$ are "close".
- $e_{\psi}\left(N_{\varphi}(G)\right)<M^{2} \cdot 2^{-n / 100}$ unless $(\varphi, \psi)$ is "exceptional".
- Why do we need the first case? $\Rightarrow\left|S_{n}\right|=n!\gg 2^{n}$
- Exceptional: $\psi^{2}$ is identity, $\varphi \approx \psi, \ldots$
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$\Rightarrow \exists \psi$-clique of size at least $e_{\psi}\left(N_{\varphi}(G)\right) /\left|N_{\varphi}(G)\right|$.


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## Lemma (without proof)

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( $\alpha$ ) Either $\left|N_{\varphi}(G)\right|<M \cdot 2^{-n / 200}$ for all $G \in \mathcal{G}, \varphi \in S_{n}$;
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- In particular, $\psi$ is an involution.
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- Assume $(\alpha)$ happens $((\beta)$ is more complicated).


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Condition: $\left|N_{\varphi}(G)\right|<M \cdot 2^{-n / 200}$ for all $G \in \mathcal{G}, \varphi \in S_{n}$.

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- For other $\psi, \#\left\{G_{1} \stackrel{\psi}{\cong} G_{2}\right\} \leq\left|\mathcal{G} \backslash N_{\varphi}\left(G_{0}\right)\right| \cdot M \cdot n^{-10 n}$.


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Condition: $M=2^{\left.\frac{1}{2}\binom{n}{2}-\frac{n}{2}\right)},|\mathcal{G}| / n!\leq\left|N_{\varphi}\left(G_{0}\right)\right|<M \cdot 2^{-n / 100}$.


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\end{aligned}
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## Further directions

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## Any suggestions?

## The End

## Questions? Comments?

